

Hybrid-relativism and revenge

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Outline

I. Absolute generality – an introduction

II. Relativist-friendly modality

III. Hybrid relativism – regimenting the argument from IE

IV. A revenge problem for hybrid relativism

Absolute generality

Quantifiers are often restricted:

- (1) No donkey talks.
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- (3) Everything is mereologically simple.
 - (4) Everything belongs to a set.
- \forall -absolutist: some domain comprises absolutely everything
 - \forall -relativist: no domain comprises absolutely everything

Relativism – the case from indefinite extensibility (IE)

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The Russell Reductio

(Zermelo 1908)

Given a domain M , let $r = \{x \in M : x \notin x\}$. Then $r \notin M$

(no assumptions about M)

Orthodox absolutist response – no Russell set r

- suppose M is the intended domain of impure set theory ZFCU
- not every plurality in M forms a set (M not a set)

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Orthodox relativist response – no absolutely comprehensive M

- sequence of intended models of ZFCU with ever-larger domains:
 $M_0, M_1, M_2 \dots$
- any plurality in M_α ‘eventually’ forms a set

How to regiment \forall -relativism?

Relativist (try 1): no set-domain is comprehensive

$$\neg \exists s \forall x (x \in s)$$

- harmless: absolutists deny that domains must be set-domains

How to regiment \forall -relativism?

Relativist (try 1): no set-domain is comprehensive $\neg\exists s\forall x(x \in s)$

- harmless: absolutists deny that domains must be set-domains
- alternative: domains are ‘plurality’-domains.

Absolutist: some plurality-domain is comprehensive

CD: $\exists xx\forall x(x < xx)$ ($\exists xx$ ‘zero or more’; $<$ ‘one of’)

Relativist (try 2): no plurality-domain is comprehensive (i.e. \neg CD)

- hopeless: standard plural logic – e.g. PFO – refutes \neg CD
- CD is a consequence of an axiom of PFO – namely, PC:

Plural Comprehension: any $\phi(x)$ defines a ‘plurality’

PC: $\exists xx(xx \equiv_x \phi(x))$ ($xx \equiv_x \phi(x) := \forall x(x < xx \leftrightarrow \phi(x))$)

- response: frame relativism in a modal language

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Relativist-friendly modality

The plural language of ZFCU: \mathcal{L}

- singular and plural quantifiers $\forall v$, $\forall vv$
- non-logical predicates: β and \in

MT-structure: $\mathcal{M} := \langle M, S, E \rangle$ $(M \neq \emptyset, S \subseteq M, E \subseteq M \times M)$

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A modal extension: \mathcal{L}^\diamond

- adds ‘forward-’ and ‘backwards-looking’ modal operators
 - $\Box_{>} \psi :=$ ‘ ψ will always hold at later stages/models’
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MT-hierarchy: $\mathcal{H} := \{\mathcal{M}_\alpha\}_{\alpha < \lambda}$ $(\lambda \text{ is a limit ordinal})$

- \mathcal{M}_α is an MT-structure or $\langle \emptyset, \emptyset, \emptyset \rangle$ $(M_\alpha \neq \emptyset \text{ for some } \alpha)$
- $\alpha < \beta \Rightarrow \mathcal{M}_\alpha$ is a substructure of \mathcal{M}_β
(i.e. $M_\alpha \subseteq M_\beta, E_\beta \cap M_\alpha = E_\alpha$, etc.)

Kripke semantics

Given $\mathcal{H} := \{M_\alpha\}_{\alpha < \lambda}$, $\alpha < \lambda$, suitable σ :

- $u \in v$ is $\text{true}_{\alpha, \sigma}$ iff $\langle \sigma(u), \sigma(v) \rangle \in E_\alpha$ (similarly for β , $<$, $=$)
 - $\forall v \psi$ is $\text{true}_{\alpha, \sigma}$ iff $\forall a \in M_\alpha$: ψ is $\text{true}_{\alpha, \sigma[v/a]}$
 - $\forall v \forall \psi$ is $\text{true}_{\alpha, \sigma}$ iff $\forall A \subseteq M_\alpha$: ψ is $\text{true}_{\alpha, \sigma[vv/A]}$
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- $\Box_{>} \psi$ is $\text{true}_{\alpha, \sigma}$ iff $\forall \beta > \alpha$: ψ is $\text{true}_{\beta, \sigma}$ ($\beta < \lambda$)
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- disclaimers:
 - \Box not semantically reducible to $\forall \beta$.
 - ‘non-circumstantial’ modality (despite ‘process’-metaphor)

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- disclaimers:
 - \Box not semantically reducible to $\forall \beta$.
 - ‘non-circumstantial’ modality (despite ‘process’-metaphor)
- target: ‘hybrid’ relativists, who accept \Box (or $\Box_{>}$) as primitive

A logic for \mathcal{L}^\diamond

Bimodal logic – MPFO

- negative, free PFO (with PC) + basic tense logic (defined $E!v$ and $E!vv$)

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Bimodal logic – MPFO

- negative, free PFO (with PC) + basic tense logic (defined $E!v$ and $E!vv$)

- D: $\Box_{>} \psi \rightarrow \Diamond_{>} \psi$ ($\Diamond_{>} \psi := \neg \Box_{>} \neg \psi$)
- H: $M\psi_1 \wedge M\psi_2 \rightarrow (M(\psi_1 \wedge \psi_2) \vee M(\psi_1 \wedge M\psi_2) \vee M(\psi_2 \wedge M\psi_1))$ ($M = \Diamond_{>}, \Diamond_{<}$)
- LÖB: $\Diamond_{<} \psi \rightarrow \Diamond_{<}(\psi \wedge \Box_{<} \neg \psi)$

- CBF: $\Box_{>} \forall v \psi \rightarrow \forall v \Box_{>} \psi$ (v singular or plural)
- STA- \in : $\Box \forall x \forall y (\Diamond x \in y \rightarrow x \in y)$ (also for $\beta, =, <$)

- E_1 : $\Diamond E!v$

Derived unimodal axioms

- \Box conforms to s5

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Can \forall -relativists achieve absolute generality?

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- hybrid relativist – yes, by ‘modalization’

$$[\cdot]^\diamond : \mathcal{L} \rightarrow \mathcal{L}^\diamond \quad \forall \mapsto \Box \forall \quad \exists \mapsto \Diamond \exists \quad \Phi \mapsto \Diamond \Phi \quad (\text{atomic } \Phi)$$

(4) Everything belongs to a set. absolutely general?

(4f) $\forall x \exists s (x \in s)$ no – if we quantify over M_α

(4f) $^\diamond$ $\Box \forall x \Diamond \exists s (x \in^\diamond s)$ yes

(infix convention: $(x \in^\diamond s) := (x \in s)^\diamond$)

Mirroring (first-order): $\Gamma \vdash_{\text{FOL}} \phi$ iff $\{\gamma^\diamond : \gamma \in \Gamma\} \vdash_{\text{MPFO}} \phi^\diamond$

- hybrid relativism = \forall -relativism + $\Box\forall$ -absolutism (‘HRism’)

- \forall -relativism: \forall never has an absolutely comprehensive domain
- $\Box\forall$ -absolutism: $\Box\forall$ ‘generalizes’ about absolutely everything

How to regiment \forall -relativism? – revisited

Recall the second attempt:

Relativism (try 2): no plurality-domain comprises everything

$\neg_{\text{CD}} : \quad \neg \exists x x \forall x (x < x x)$

- MPFO – with PC – refutes try 2

The argument from IE – revisited

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Naive Comprehension: every plurality is collected.

$$(NCA) \quad \forall xx \exists s (s \equiv xx) \quad (s \equiv xx := \forall x (x \in s \leftrightarrow x < xx))$$

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HRist may instead deploy a modal plenitude principle:

Naive Comprehension \diamond : absolutely every plurality is collectable.

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- $\text{NCA}^\diamond \not\vdash_{\text{MPFO}} \perp$

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What makes uncollectable xx uncollectable?

Consider e.g. the plurality oo of ordinals (in the domain M)

Answer 1 – the paradoxes: e.g. $ZFCU \vdash \neg \exists s \forall x (x \in s \leftrightarrow \text{ORD}(x))$

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Relativist objects – ZFCU tells us about collectedness not collectability

- oo is not collected $(\text{ZFCU} \vdash_{\text{PFO}} \neg \exists s (s \equiv oo))$
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Relativist objects – why that order-type?

- it's coherent for oo to form a set (at a stage outside \mathcal{S})
- then there'd be a longer sequence \mathcal{S}^+
- why privilege \mathcal{S} over \mathcal{S}^+ ?

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- question: when is $\phi^\diamond(x)$ plurally comprehensible?
- answer: precisely when $\phi^\diamond(x)$ is not indefinitely extensible

- $\text{EXT}_x[\phi^\diamond(x)] := \diamond_{>} \exists x(\phi^\diamond(x) \wedge \square_{<} \neg E!x) \quad (E!x := x = x)$

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Hybrid-relativism and revenge

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- do similar considerations point to r outside $\Box \forall x$?

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TGRist – revenge problem for HR

- introduce \Box , intended to be broader than \square
- frame a new version of argument from IE
- argue that non-comprehensibility isn't much easier to explain than non-collectability

More relativist-friendly modality

An extended modal language – \mathcal{L}^\diamond

- add $\Box_{>}$ and $\Box_{<}$ to \mathcal{L}^{\diamond} (\Box now sole primitive)
- define $\Box\psi := \Box_{>}\psi \wedge \psi \wedge \Box_{<}\psi$

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- add $\Box_>$ and $\Box_<$ to \mathcal{L}^\diamond (\Box now sole primitive)
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Extended MT-hierarchy: $\mathcal{H}_{\lambda,\lambda^+} := \langle \mathcal{H}, \lambda \rangle$

- $\mathcal{H}_\lambda := \{\mathcal{M}_\alpha\}_{\alpha < \lambda}$ is an MT-hierarchy
- $\mathcal{H}_{\lambda^+} := \mathcal{H} = \{\mathcal{M}_\alpha\}_{\alpha < \lambda^+}$ is an MT-hierarchy
- $\lambda \leq \lambda^+$

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Given $\mathcal{H}_{\lambda,\lambda^+}$, $\alpha < \lambda^+$, suitable σ :

- $\Box\psi$ is $\text{true}_{\alpha,\sigma}$ iff $\forall\beta: \psi$ is $\text{true}_{\beta,\sigma}$ ($\beta < \lambda^+$)
- similarly for $\Box_>$ and $\Box_<$
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A logic for \mathcal{L}^\diamond

Trimodal logic – MPFO⁺

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- $X_b: \Box\psi \rightarrow \Box\psi$

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- MPFO for $\Box_>$ and $\Box_<$.

- X_b : $\Box\psi \rightarrow \Box\psi$

- D: $\Box\psi \rightarrow \Diamond\psi$

($\Box\psi \rightarrow \psi$ may fail above λ)

- 4_b : $\Box\psi \rightarrow \Box\Box\psi$

- 5_b : $\Diamond\psi \rightarrow \Box\Diamond\psi$

- $\text{STA}^+ \text{-}\epsilon$: $\Box\forall u\Box\forall v(\Diamond u \in v \rightarrow \Diamond u \in v)$

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- $STA^+-\epsilon$: $\Box\forall u\Box\forall v(\Diamond u \in v \rightarrow \Diamond u \in v)$ (also for $\beta, <, =$)

- E_1^- : $\Box\forall v\Diamond E!v$ ($\Diamond E!v$ may fail)

A revenge argument

- TGRist conclusion:

New Items: not \Box -absolutely every item is a \Diamond -potential item

NI: $\neg \Box \forall x \Diamond E!x$

- NI follows (in MPFO^+) from two premisses:

Naive Comprehension $^\Diamond$ – **NCA $^\Diamond$** : $\Box \forall xx \Diamond \exists s (s \equiv^\Diamond xx)$

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Liberal Comprehensibility: each $+$ -invariant condition $\psi(x)$ that is not $+$ -extensible is $+$ -comprehensible.

LC: $+INV[\psi(x)] \wedge \neg +EXT_x[\psi(x)] \rightarrow \Diamond \exists xx (xx \equiv_x^\Diamond \psi(x))$

$+INV[\psi] := \Box \forall \vec{v} (\Box \psi \vee \Box \neg \psi)$ $+EXT_x[\psi(x)] := \Diamond \exists x (\psi(x) \wedge \Box \neg E!x)$

- HRist response: deny LC

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- HRist response: deny LC – explanatory challenge ...

What makes non-comprehensible $\phi^\diamond(x)$ non-comprehensible?

Answer (Linnebo): some conditions have ‘an inexhaustible character that renders them intrinsically unsuitable for defining sets.’

(‘Potential Hierarchy of Sets’, p. 213)

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Inextensibility: $\text{INV}[\psi(x)] \vdash \Box \text{EXT}_x[\psi(x)] \rightarrow \neg \Diamond \exists xx(xx \equiv_r^\Diamond \psi(x))$

- consider $\text{ORD}^\diamond(x)$. $(\vdash \text{ORD}^\diamond(x) \rightarrow \diamond \beta x)$
- ‘inexhaustible character’ – or, at least $\Box \text{EXT}_x[\text{ORD}^\diamond(x)]$
- so, $\text{ORD}^\diamond(x)$ is non- \diamond -comprehensible $(\neg \diamond \exists x x(x x \equiv_x^\diamond \text{ORD}^\diamond(x)))$

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- consider $\text{ORD}^\diamond(x)$. $(\vdash \text{ORD}^\diamond(x) \rightarrow \diamond \beta x)$

- ‘inexhaustible character’ – or, at least $\Box \text{EXT}_x[\text{ORD}^\diamond(x)]$
- so, $\text{ORD}^\diamond(x)$ is non- \diamond -comprehensible $(\neg \diamond \exists x x (xx \equiv_x^\diamond \text{ORD}^\diamond(x)))$

But is $\text{ORD}^\diamond(x)$ +-comprehensible? (i.e. $\diamond \exists x x (xx \equiv_x^\diamond \text{ORD}^\diamond(x))$)

What makes non-comprehensible $\phi^\diamond(x)$ non-comprehensible?

Answer (Linnebo): some conditions have ‘an inexhaustible character that renders them intrinsically unsuitable for defining sets.’

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- TGRist: yes, since $\neg +\text{EXT}_x[\text{ORD}^\diamond(x)]$ (so $\neg \Box \text{EXT}_x[\text{ORD}^\diamond(x)]$)

How convincing is this HRist explanation?

What makes $\text{ORD}^\diamond(x)$ \square -indefinitely extensible? (s.t. $\square_{\text{EXT}_x}[\text{ORD}^\diamond(x)]$)

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- relativist objects – why that hierarchy?
 - it's coherent for $\text{ORD}^\diamond(x)$ to be +-comprehended (outside \mathcal{H})
 - then there'd be a more inclusive hierarchy \mathcal{H}^+
 - why privilege \mathcal{H} over \mathcal{H}^+ ?